

# **CFRG Specifications**

## **In theory and practice**

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**Disclaimer: This is based on my personal experience accumulated in the CFRG, and does not represent the group's shared view**

**Goal: Highlight ways we can improve the group's primary deliverables, not point fingers or assign blame**

**What is the problem?**

# CFRG Specifications

## In theory

### CFRG charter

The CFRG serves as a bridge between theory and practice, bringing new cryptographic techniques to the Internet community and promoting an understanding of the use and applicability of these mechanisms via **Informational RFCs**...

### RFC2026

An "**Informational**" specification is published for the **general information of the Internet community**, and does not represent an Internet community consensus or recommendation. The Informational designation is ... **subject only to editorial considerations and to verification that there has been adequate coordination with the standards process.**

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# CFRG Specifications

## In practice

Specifications have significant impact on protocol design, security analysis, and implementations:

RFC2104: HMAC

RFC5869: HKDF

RFC7748: Curve25519/X25519

RFC8032: EdDSA

RFC9180: HPKE

Specifications target a wide variety of audiences:

Protocol designers and implementers

Cryptographic reviewers

...

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# CFRG Specifications

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Specifications target a wide variety of audiences

High stakes

Quality is critical

RFC2104:

RFC5869:

RFC7748:

RFC8032: EdDSA

RFC9180: HPKE

...

implementers



**Specification quality**

# Specification Components

There are (at least) three different parts of a specification:

1. Functional specification. What does this object do? What is its purpose?
2. Syntax specification. How do I interact with this object?
3. Implementation specification. How does this object work internally?

Presentation of each should be tailored to its audience

# Guiding Questions

1. Is the specification **easy to understand and use**?

Is the functional description of the cryptographic object clear?

What is the cognitive load required to understand the specification?

Is the syntax of the object clear?

2. Will the specification yield **consistent and correct implementations**?

Is the functional behavior well-defined?

Is the syntax correct and amenable to type-safe implementations?

Is the implementation description clear?

# Example: RFC8032 (EdDSA)

## 5.1.7. Verify

1. To verify a signature on a message  $M$  using public key  $A$ , with  $F$  being 0 for Ed25519ctx, 1 for Ed25519ph, and if Ed25519ctx or Ed25519ph is being used,  $C$  being the context, first split the signature into two 32-octet halves. Decode the first half as a point  $R$ , and the second half as an integer  $S$ , in the range  $0 \leq s < L$ . Decode the public key  $A$  as point  $A'$ . If any of the decodings fail (including  $S$  being out of range), the signature is invalid.
2. Compute  $\text{SHA512}(\text{dom2}(F, C) \parallel R \parallel A \parallel \text{PH}(M))$ , and interpret the 64-octet digest as a little-endian integer  $k$ .
3. Check the group equation  $[8][S]B = [8]R + [8][k]A'$ . It's sufficient, but not required, to instead check  $[S]B = R + [k]A'$ .

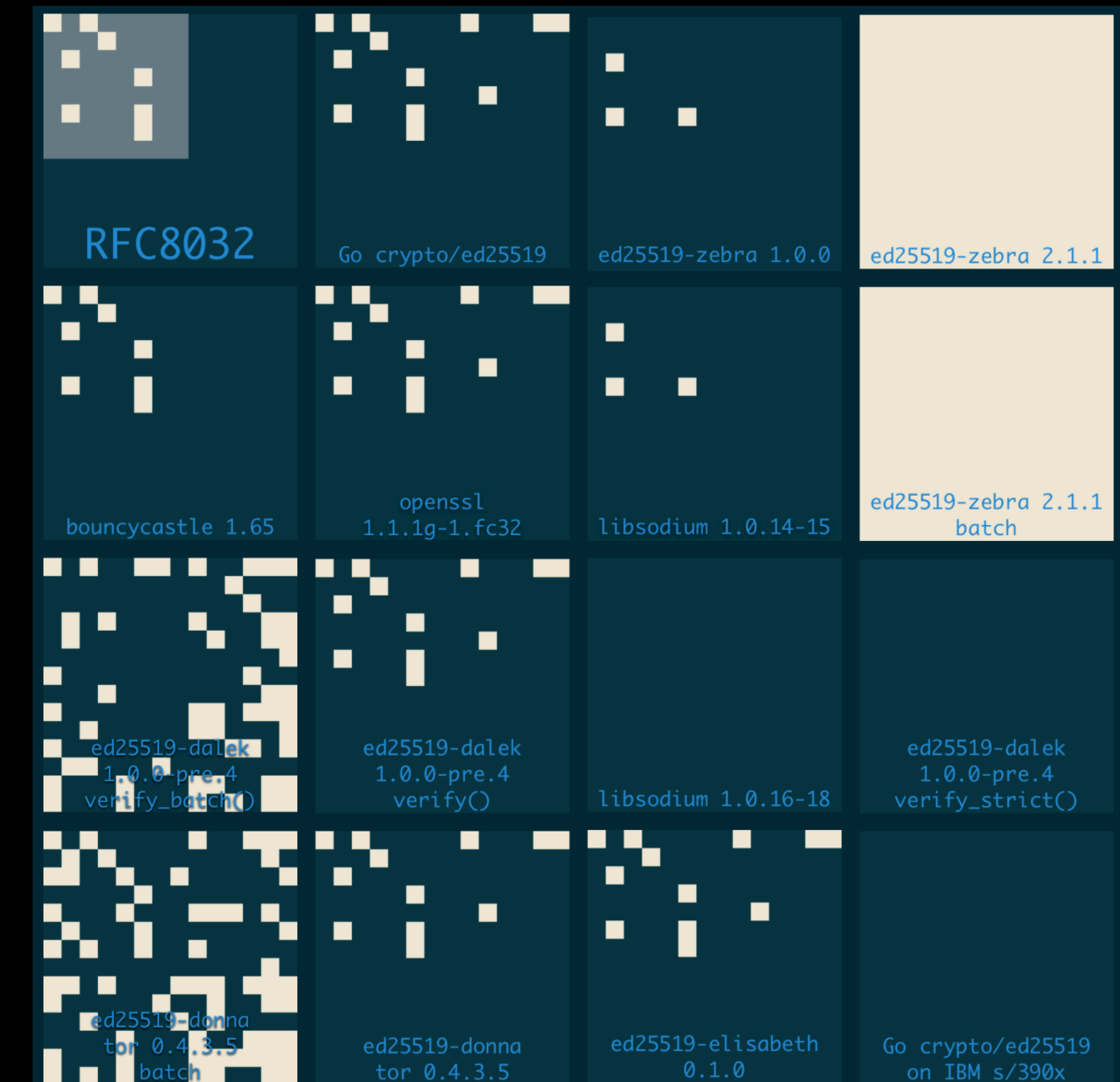
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Source: <https://hdevalence.ca/blog/2020-10-04-its-25519am>

1. Is the specification easy to understand and use? ✓

2. Will the specification yield consistent and correct implementations? ✗

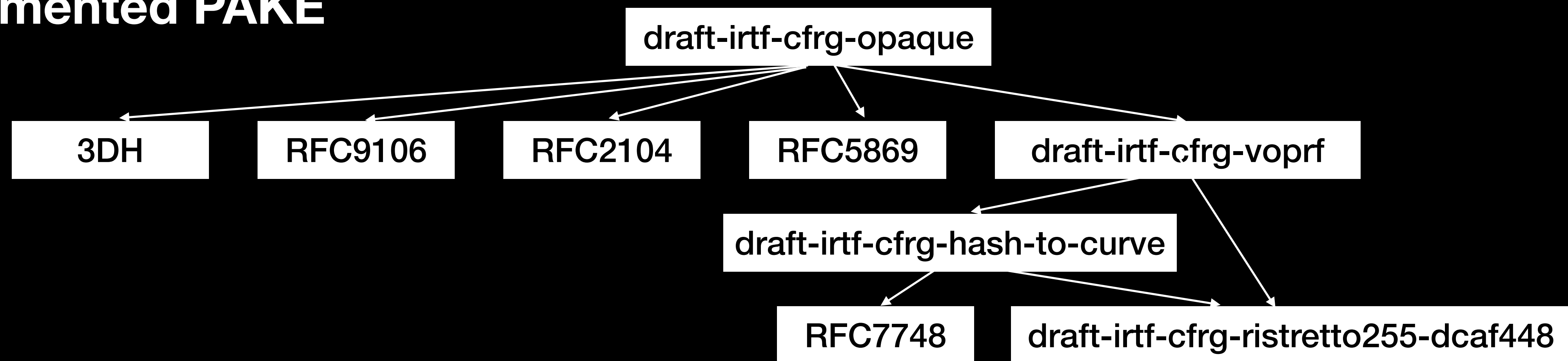
# Example: OPAQUE (draft-irtf-cfrg-opaque)

OPAQUE is a compiler for translating an OPRF, hash function, memory hard function (MHF), and authenticated key exchange (AKE) protocol into a **strong, augmented PAKE**

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# Core Problem

CFRG produces technical specifications for cryptographic objects that are consumed by a diverse audience

Each object is expected to have easy-to-understand and well-defined behavior with clear syntax (API)

Failure to establish this clarity and consistency will yield specifications with little or no value and possibly harmful consequences in practice

Writing technical specifications such that they detail cryptographic objects with **well-defined behavior** and **clear syntax** is challenging



# Case study: hash-to-curve

# Hash-to-Curve Overview

`hash_to_curve` is a uniform encoding from byte strings to points in some elliptic curve group  $G$

```
hash_to_curve(msg)
```

```
Input: msg, an arbitrary-length byte string.
```

```
Output: P, a point in G.
```

```
Steps:
```

```
1. u = hash_to_field(msg, 2)
```

```
2. Q0 = map_to_curve(u[0])
```

```
3. Q1 = map_to_curve(u[1])
```

```
4. R = Q0 + Q1 # Point addition
```

```
5. P = clear_cofactor(R)
```

```
6. return P
```

1. Is the specification easy to understand and use?
2. Will the specification yield consistent and correct implementations?

# For the specification user...

## Functional description

```
hash_to_curve(msg)
```

Input: msg, an arbitrary-length byte string.

Output: P, a point in G.

Steps:

1.  $u = \text{hash\_to\_field}(\text{msg}, 2)$
2.  $Q_0 = \text{map\_to\_curve}(u[0])$
3.  $Q_1 = \text{map\_to\_curve}(u[1])$
4.  $R = Q_0 + Q_1$  # Point addition
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# For the specification user...

## Functional description

`hash_to_curve(msg)`

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Steps:

1. `u = hash_to_field(msg, 2)`
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3. `Q1 = map_to_curve(u[1])`
4. `R = Q0 + Q1` # Point addition
5. `P = clear_cofactor(R)`
6. return `P`

`hash_to_field(msg, count)`

Inputs:

- `msg`, a byte string containing the message to hash.
- `count`, the number of elements of  $F$  to output.

Outputs:

- `(u_0, ..., u_(count - 1))`, a list of field elements.

Steps: defined in [Section 5](#).

`map_to_curve(u)`

Input: `u`, an element of field  $F$ .

Output: `Q`, a point on the elliptic curve  $E$ .

Steps: defined in [Section 6](#).

`clear_cofactor(Q)`

Input: `Q`, a point on the elliptic curve  $E$ .

Output: `P`, a point in  $G$ .

Steps: defined in [Section 7](#).

Functional specification should be maximally clear for people trying to understand what the object does without understanding how it works

# For the API designer...

## Syntax specification

```
hash_to_curve(msg)
```

Input: msg, an arbitrary-length byte string.

Output: P, a point in G.

Steps:

1. `u = hash_to_field(msg, 2)`
2. `Q0 = map_to_curve(u[0])`
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4. `R = Q0 + Q1` # Point addition
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### 6.7.1. Elligator 2 method

Bernstein, Hamburg, Krasnova, and Lange give a mapping that applies to any curve with a point of order 2 [BHKL13], which they call Elligator 2.

Preconditions: A Montgomery curve  $K * t^2 = s^3 + J * s^2 + s$  where  $J \neq 0$ ,  $K \neq 0$ , and  $(J^2 - 4) / K^2$  is non-zero and non-square in  $F$ .

Constants:

- \* J and K, the parameters of the elliptic curve.
- \* Z, a non-square element of  $F$ . [Appendix H.3](#) gives a Sage [SAGE] script that outputs the RECOMMENDED Z.

Sign of t: this mapping fixes the sign of t as specified in [BHKL13]. No additional adjustment is required.

Exceptions: The exceptional case is  $Z * u^2 == -1$ , i.e.,  $1 + Z * u^2 == 0$ . Implementations must detect this case and set  $x1 = -(J / K)$ . Note that this can only happen when  $q = 3 \pmod{4}$ .

Syntax specification should follow from the functional specification and be easy to use and hard to misuse

# For the implementer...

## Implementation description

hash\_to\_curve(msg)

Input: msg, an arbitrary-length byte string.

Output: P, a point in G.

Steps:

1.  $u = \text{hash\_to\_field}(\text{msg}, 2)$
2.  $Q_0 = \text{map\_to\_curve}(u[0])$
3.  $Q_1 = \text{map\_to\_curve}(u[1])$
4.  $R = Q_0 + Q_1$  # Point addition
5.  $P = \text{clear\_cofactor}(R)$
6. return P

### F.3. Elligator 2 method

This section gives a straight-line implementation of the Elligator 2 method for any Montgomery curve of the form given in [Section 6.7](#). See [Section 6.7.1](#) for information on the constants used in this mapping.

[Appendix G.2](#) gives optimized straight-line procedures that apply to specific classes of curves and base fields, including curve25519 and curve448 [[RFC7748](#)].

map\_to\_curve\_elligator2(u)

Input: u, an element of F.

Output: (s, t), a point on M.

Constants:

1.  $c_1 = J / K$
2.  $c_2 = 1 / K^2$

Steps:

1.  $tv_1 = u^2$
2.  $tv_1 = Z * tv_1$  #  $Z * u^2$
3.  $e_1 = tv_1 == -1$  # exceptional case:  $Z * u^2 == -1$
4.  $tv_1 = \text{CMOV}(tv_1, 0, e_1)$  # if  $tv_1 == -1$ , set  $tv_1 = 0$
5.  $x_1 = tv_1 + 1$
6.  $x_1 = \text{inv0}(x_1)$
7.  $x_1 = -c_1 * x_1$  #  $x_1 = -(J / K) / (1 + Z * u^2)$
8.  $gx_1 = x_1 + c_1$
9.  $gx_1 = gx_1 * x_1$
10.  $gx_1 = gx_1 + c_2$
11.  $gx_1 = gx_1 * x_1$  #  $gx_1 = x_1^3 + (J / K) * x_1^2 + x_1 / K^2$
12.  $x_2 = -x_1 - c_1$
13.  $gx_2 = tv_1 * gx_1$

Implementation specification should given implementers confidence in their implementation



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*A way forward*



# Remember the Audience

## Align presentation styles

Aim towards alignment between across functional, syntax, and implementation specifications

- Use consistent pseudocode to describe all three

- Make pseudocode match reference implementation as close as possible

# Consistency is Key

## Reduce cognitive load

Improve clarity by reusing concepts and notation

- Use consistent terminology and vocabulary

- Adopt consistent presentation format (e.g., for pseudocode)

CFRG strives to produce high quality specifications of cryptographic objects

... but consistency across drafts is lacking

# Embrace Formality

## Adopt formally verified toolchains

Minimize inconsistencies between functional and implementation descriptions

Can formally verified reference implementations (hacspec) offer a way forward?

What is the simplest, most approachable reference implementation format?  
(Python, C/C++, Sage, Rust, etc)

... Unclear if these existing languages with formally verified toolchains are ready for production use in long-lived specifications

# Wrapping up

# Summary

CFRG specifications (... **standards** ...) are important to the community

More work to be done to improve consistency across drafts

Share terminology, concepts, and notation? Share reference implementations?

Reference implementation requirements and reviews? Common requirements for syntax (APIs)?

Explore applicability of formal methods for specification